

# Modeling and Algorithm for Dynamic Multi-Objective Max-CSPs

Tenda Okimoto <sup>1,2</sup>, Tony Rebeiro <sup>3</sup>, Maxime Clement <sup>4</sup>  
and Katsumi Inoue <sup>2</sup>

<sup>1</sup> Transdisciplinary Research Integration Center

<sup>2</sup> National Institute of Informatics

<sup>3</sup> The Graduate University for Advanced Studies

<sup>4</sup> Universite Pierre et Marie Curie

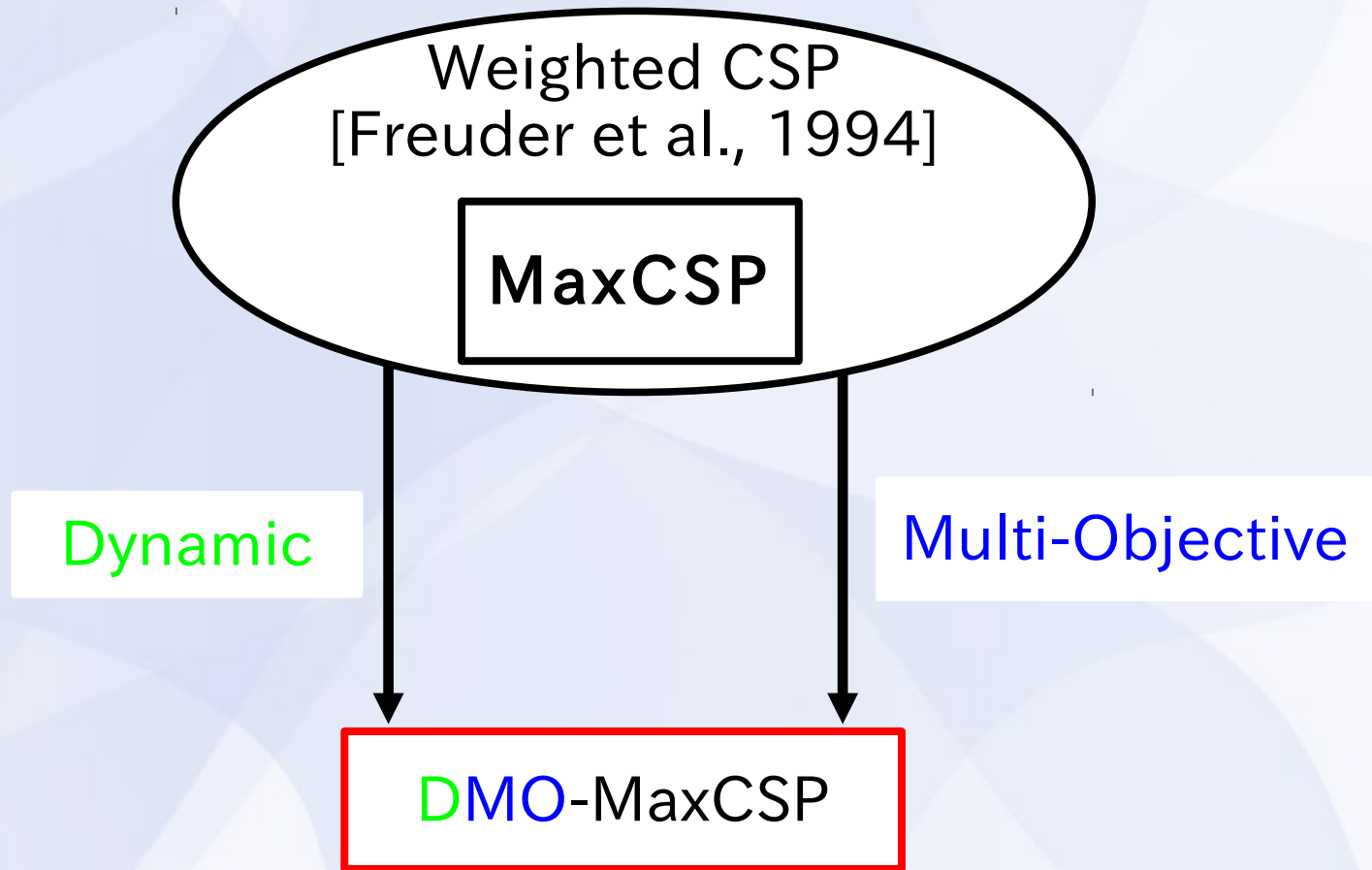
# On going

- Complete/Incomplete algorithms for MO-DCOPs
  - Proof for generalisation
  - Comparison is required with B-MOMS
- AOF-technique based Algorithm for Dynamic MO-DCOPs
  - Diversity : method to provide weights
- Modeling and Algorithm for Dynamic MO-DCOPs
  - Reactive approach : everything can be changed (s-robustness)
- Modeling and Algorithm for Dynamic MO-Max-CSPs
  - Proactive approach: with Nicolas
  - Reactive approach: with Tony
- Prism problem (extension of WSR-paper with Ikegai)

# Motivation

- Many real world problems involve multiple criteria that should be considered separately.
- Many real systems are dynamic, i.e., problems change at runtime.

# Related Works



# Summary

- Provide a model for DMO-MaxCSPs.
  - Defined by a sequence of MO-MaxCSPs.
  - # variables / domains / # constraints / # objectives can be changed.

# Summary

- Provide a model for DMO-MaxCSPs.
  - Defined by a sequence of MO-MaxCSPs.
  - # variables / domains / # constraints / # objectives can be changed.
- Develop the first algorithm for DMO-MaxCSPs.
  - Find a sequence of a set of solutions based on new solution criteria l-, m- and s-robustness.
  - The complexity is bounded by the parameter s.

# Outline

- Preliminaries
  - CSP/Max-CSP
- Multi-Objective Max-CSP :
  - $(l,m)$ -robustness
- Dynamic MO-MaxCSP (DMO-MaxCSP)
  - $s$ -weak robustness/ $(l,m,s)$ -robustness
- Algorithm for DMO-MaxCSPs
- Experiments
- Conclusion and future work

# Outline

- Preliminaries
  - CSP/Max-CSP
- Multi-Objective Max-CSP :
  - $(l,m)$ -robustness
- Dynamic MO-MaxCSP (DMO-MaxCSP)
  - $s$ -weak robustness/ $(l,m,s)$ -robustness
- Algorithm for DMO-MaxCSPs
- Experiments
- Conclusion and future work



# CSP/Max-CSP

- **CSP/Max-CSP** is defined by  $\langle V, D, C \rangle$  where
  - $V$  : a set of variables,
  - $D$  : a set of domains,
  - $C$  : a set of constraints.
- **Goal for CSP**: Find a consistent assignment of values to variables.

# CSP/Max-CSP

- **CSP/Max-CSP** is defined by  $\langle V, D, C \rangle$  where
  - $V$  : a set of variables,
  - $D$  : a set of domains,
  - $C$  : a set of constraints.
- **Goal for CSP**: Find a consistent assignment of values to variables.
- **Goal for Max-CSP** : Find an assignment which satisfies as many constraints as possible.
  - Find an assignment which minimizes the number of violated constraints.

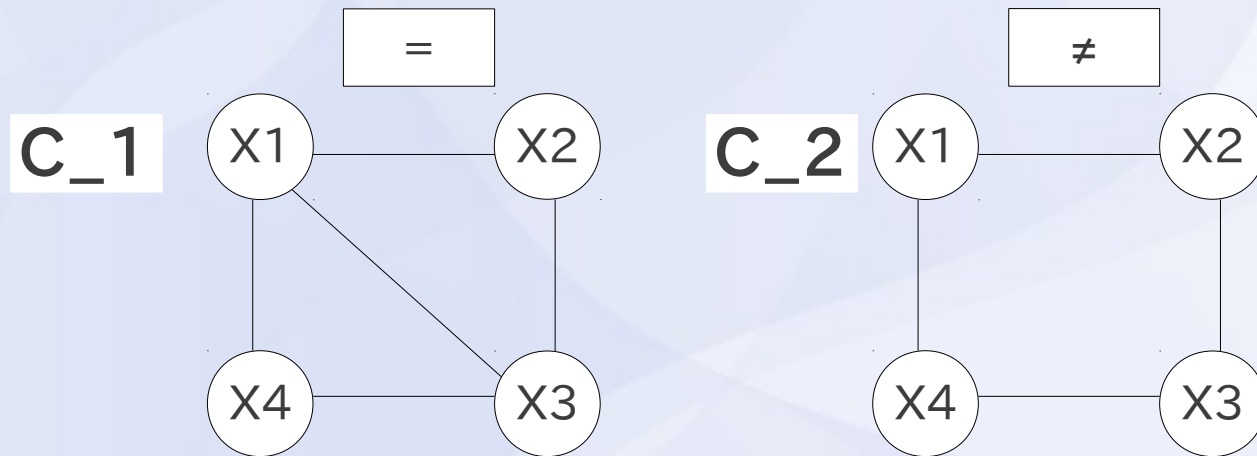
# Outline

- Preliminaries
  - CSP/Max-CSP
- **Multi-Objective Max-CSP :**
  - **(l,m)-robustness**
- Dynamic MO-MaxCSP (DMO-MaxCSP)
  - s-weak robustness/(l,m,s)-robustness
- Algorithm for DMO-MaxCSPs
- Experiments
- Conclusion and future work

# Multi-Objective Max-CSP (Formalisation)

- MO-MaxCSP is defined by  $\langle V, D, \Gamma \rangle$  where
  - $V$  : a set of variables,
  - $D$  : a set of domains,
  - $\Gamma$  : a set of sets of constraints  $\{C_1, \dots, C_k\}$   
( $k$  criteria/objectives which cannot merge)
- Solving an MO-MaxCSP is to find a Pareto front.

# Multi-Objective Max-CSP (Example)



- Domain : {  $\bigcirc$ ,  $\bullet$  }
- $\{(x_1, \bigcirc), (x_2, \bullet), (x_3, \bigcirc), (x_4, \bullet)\} = (4, 0)$
- $\{(x_1, \bigcirc), (x_2, \bigcirc), (x_3, \bigcirc), (x_4, \bigcirc)\} = (0, 4)$

# Multi-Objective Max-CSP (Properties)

- For a cost vector  $R(A)$  and a constant  $l$ , we call that  $R(A)=(v_1, v_2, \dots, v_n)$  is a  **$l$ -weak robust solution** if each value  $v_i$  ( $1 \leq i \leq n$ ) is below  $l$ .
  - e.g.  $l=3$  : **No** (5,0) / **Yes** (2,2)

# Multi-Objective Max-CSP (Properties)

- For a cost vector  $R(A)$  and a constant  $l$ , we call that  $R(A)=(v_1, v_2, \dots, v_n)$  is a  $l$ -weak robust solution if each value  $v_i$  ( $1 \leq i \leq n$ ) is below  $l$ .
  - e.g.  $l=3$  : No  $(5,0)$  / Yes  $(2,2)$
- For a cost vector  $R(A)$  and a constant  $l$ , we call that  $R(A)=(v_1, v_2, \dots, v_n)$  is a  **$m$ -weak robust solution** if  $v_1 + v_2 + \dots + v_n$  is below  $m$ .
  - e.g.  $m=5$  : No  $(3,3)$  / Yes  $(4,1)$

# Multi-Objective Max-CSP (Properties)

- For a cost vector  $R(A)$  and a constant  $l$ , we call that  $R(A)=(v_1, v_2, \dots, v_n)$  is a  $l$ -weak robust solution if each value  $v_i$  ( $1 \leq i \leq n$ ) is below  $l$ .
  - e.g.  $l=3$  : No (5,0) / Yes (2,2)
- For a cost vector  $R(A)$  and a constant  $l$ , we call that  $R(A)=(v_1, v_2, \dots, v_n)$  is a  $m$ -weak robust solution if  $v_1 + v_2 + \dots + v_n$  is below  $m$ .
  - e.g.  $m=5$  : No (3,3) / Yes (4,1)
- $R(A)$  is a **( $l, m$ )-robust solution** if  $R(A)$  is  $l$ -weak robust solution and  $m$ -weak robust solution.
  - e.g.  $l=3, m=5$  : No (4,1) / Yes (3,2)



# Multi-Objective Max-CSP (Model)

- Given : MO-MaxCSP  $\langle V, D, \Gamma, l, m \rangle$ 
  - $V$  : a set of variables,
  - $D$  : a set of domains,
  - $\Gamma$  : a set of sets of constraints  $\{C_1, \dots, C_k\}$ ,
  - $l, m$  : constant
- Goal : Find a set of  $(l, m)$ -robust solutions.

# Outline

- Preliminaries
  - CSP/Max-CSP
- Multi-Objective Max-CSP :
  - $(l,m)$ -robustness
- **Dynamic MO-MaxCSP (DMO-MaxCSP)**
  - **$s$ -weak robustness/ $(l,m,s)$ -robustness**
- Algorithm for DMO-MaxCSPs
- Experiments
- Conclusion and future work

# Dynamic MO-MaxCSP (Formalisation)

- DMO-MaxCSP is defined by a sequence of MO-MaxCSPs.

$\langle \text{MO-MaxCSP}_1, \dots, \text{MO-MaxCSP}_p \rangle$

- # variables/domains/# constraints/# objectives can be changed at run time.
- Goal: Find a sequence of a set of Pareto front.

$\langle \text{PF}_1, \dots, \text{PF}_p \rangle$

# Dynamic MO-MaxCSP (Properties)

- Let  $R(A)$  and  $R(A')$  be two cost vectors of MO-MaxCSP and MO-MaxCSP'.
- $s = \text{dis}(A', A)$ : hamming distance between  $A$  and  $A'$  (# differences between two assignments).

# Dynamic MO-MaxCSP (Properties)

- Let  $R(A)$  and  $R(A')$  be two cost vectors of MO-MaxCSP and MO-MaxCSP'.
- $s = \text{dis}(A', A)$ : hamming distance between  $A$  and  $A'$  (# differences between two assignments).
- $R(A')$  is a **s-weak robust solution** if  $\text{dis}(A, A') \leq s$ .
  - e.g.  $S=1$  : **No**  $A = (\bigcirc, \bullet, \bigcirc, \bullet)$ ,  $A' = (\bigcirc, \bigcirc, \bigcirc, \bigcirc)$
  - $A = (\bigcirc, \bullet, \bigcirc, \bullet)$ ,  $A' = (\bigcirc, \bullet, \bigcirc, \bullet)$  :  $\text{dis}(A, A') = 0$

# Dynamic MO-MaxCSP (Properties)

- Let  $R(A)$  and  $R(A')$  be two cost vectors of MO-MaxCSP and MO-MaxCSP'.
- $s = \text{dis}(A', A)$ : hamming distance between  $A$  and  $A'$  (# differences between two assignments).
- $R(A')$  is a **s-weak robust solution** if  $\text{dis}(A, A') \leq s$ .
  - e.g.  $S=1$  : **No**  $A = (\bigcirc, \bullet, \bigcirc, \bullet)$ ,  $A' = (\bigcirc, \bigcirc, \bigcirc, \bigcirc)$
  - $A = (\bigcirc, \bullet, \bigcirc, \bullet)$ ,  $A' = (\bigcirc, \bullet, \bigcirc, \bullet)$  :  $\text{dis}(A, A') = 0$
- $R(A')$  is a **(l,m,s)-robust solution** if  $A'$  is a (l,m)-robust solution and a s-weak robust solution.

# Dynamic MO-MaxCSP (Model)

- Given : A sequence of MO-MaxCSPs  
 $\langle \text{MO-MaxCSP}_1, \dots, \text{MO-MaxCSP}_p \rangle$  where  
 $\text{MO-MaxCSP} = \langle V, D, \Gamma, l, m, s \rangle$
- Goal: Find a sequence of a set of  $(l, m, s)$ -robust solutions.

$\langle \text{RS}_1, \dots, \text{RS}_p \rangle$

# Outline

- Preliminaries
  - CSP/Max-CSP
- Multi-Objective Max-CSP :
  - $(l,m)$ -robustness
- Dynamic MO-MaxCSP (DMO-MaxCSP)
  - $s$ -weak robustness/ $(l,m,s)$ -robustness
- **Algorithm for DMO-MaxCSPs**
- Experiments
- Conclusion and future work



# Algorithm

- Based on BnB(+softAC)
  - Solve all  $(l,m,s)$ -solutions.
  - Pruning by criteria  $(l,m)$  and  $s$ .

# Complexity

- Complexity:
  - $O(|PF| \cdot n \cdot k \cdot (|D_{\max}|^s))$  ( $0 \leq s \leq d$ ), where
  - $d$  is the depth
  - $n$  is the number of variables
- Parameters :  $k$  is the number of objectives.
  - $0 \leq l \leq n(n-1)/2$
  - $0 \leq m \leq k * \{n(n-1)/2\}$
  - $0 \leq s \leq n$

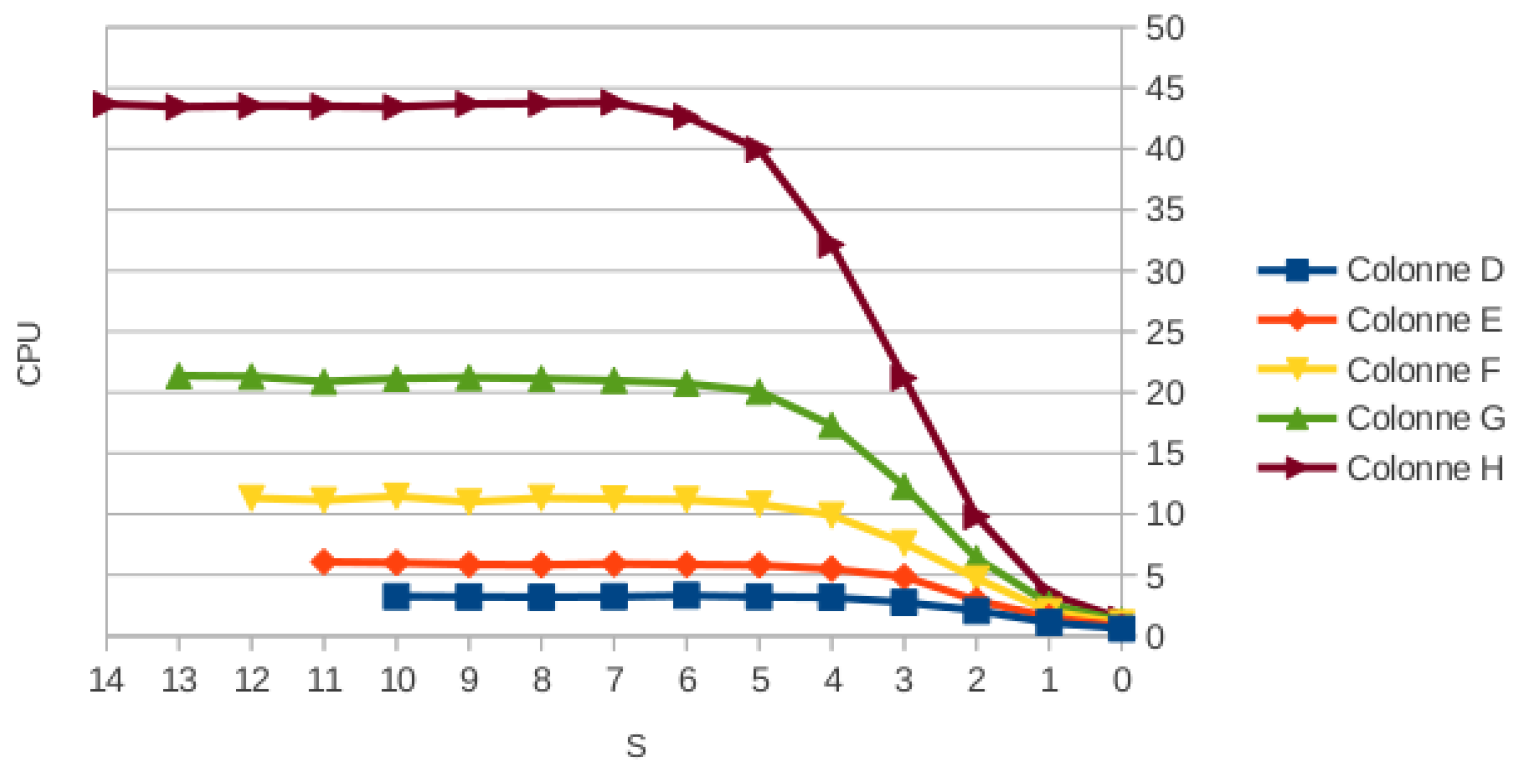
# Outline

- Preliminaries
  - CSP/Max-CSP
- Multi-Objective Max-CSP :
  - $(l,m)$ -robustness
- Dynamic MO-MaxCSP (DMO-MaxCSP)
  - $s$ -weak robustness/ $(l,m,s)$ -robustness
- Algorithm for DMO-MaxCSPs
- **Experiments**
- Conclusion and future work

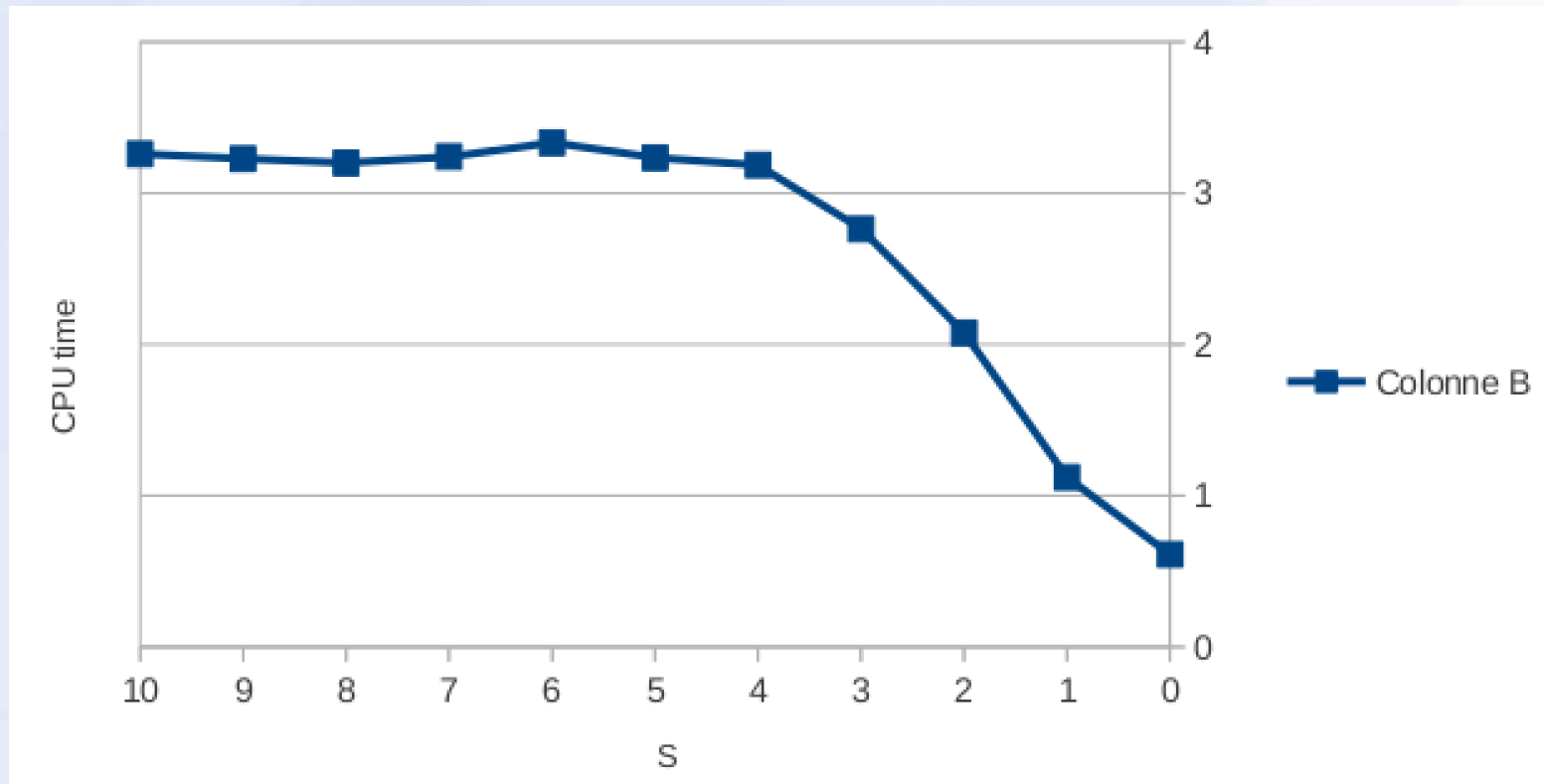
# Experiments

- Objectives : 2->3
- Domain size : 2
- Problem instances : 50
- Varying the # nodes : 10-20

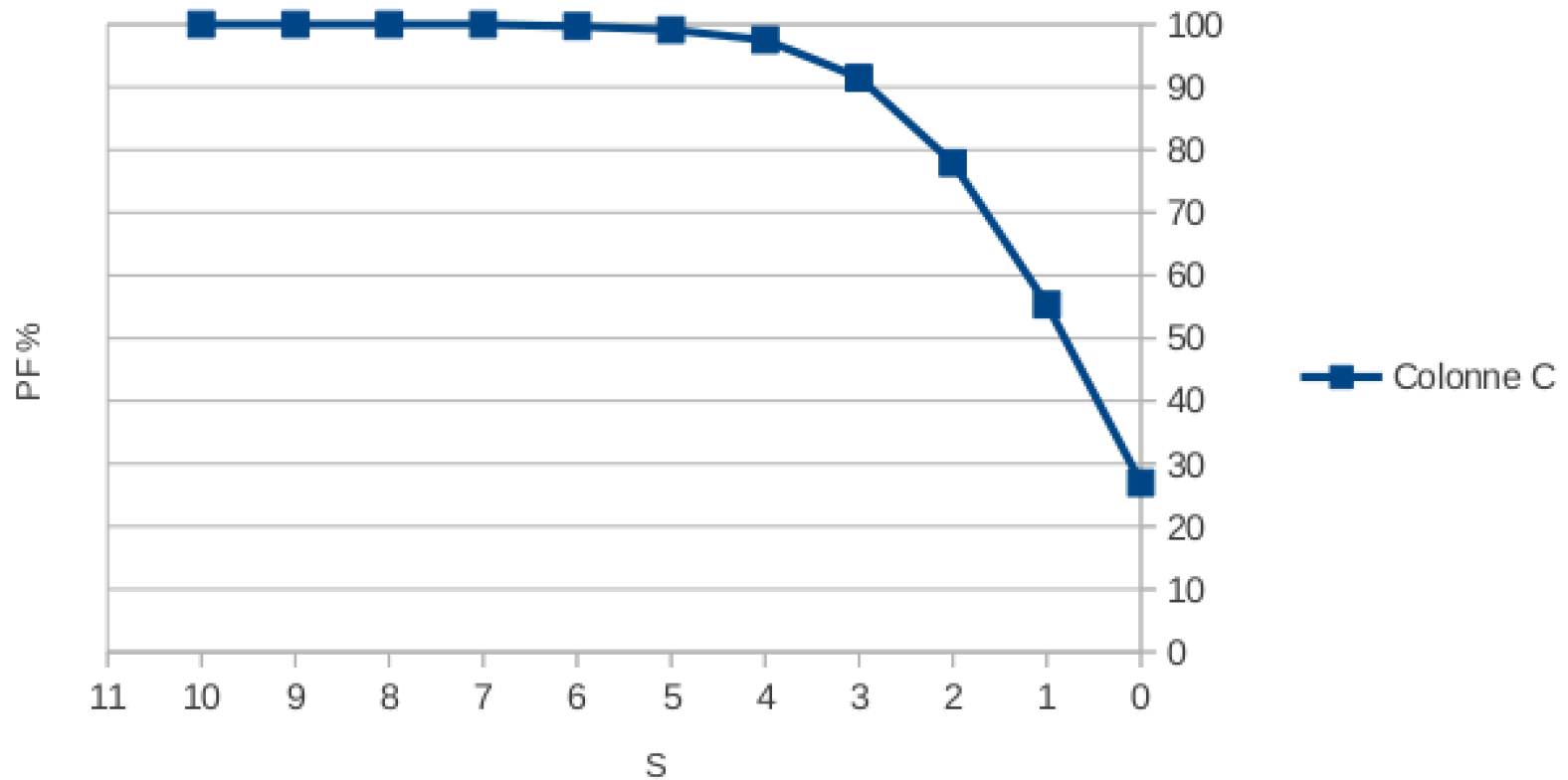
# Different # nodes



# Runtime (10 nodes)



# Pareto front (10 nodes)



# Outline

- Preliminaries
  - CSP/Max-CSP
- Multi-Objective Max-CSP :
  - $(l,m)$ -robustness
- Dynamic MO-MaxCSP (DMO-MaxCSP)
  - $s$ -weak robustness/ $(l,m,s)$ -robustness
- Algorithm for DMO-MaxCSPs
- Experiments
- **Conclusion and future work**



*Thanks!  
and  
Questions !*

*tenda@nii.ac.jp*